

Séminaire MaMux

ReactiveML

et

aspects dynamiques dans les langages synchrones

Louis Mandel

Université Paris-Sud 11
INRIA Paris-Rocquencourt

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Caractéristiques des systèmes que nous voulons programmer :

- ▶ pas de contraintes temps réel
- ▶ beaucoup de **communications et de synchronisations**
- ▶ beaucoup de **concurrency**
- ▶ **création dynamique** de processus

ReactiveML

Extension d'un langage généraliste (Ocaml *)

- ▶ structures de données
- ▶ structures de contrôle

Modèle de concurrence simple et déterministe

- ▶ composition parallèle
- ▶ communications entre processus

Compilé vers du code Ocaml

- ▶ générateur de bytecode et de code natif
- ▶ exécutif efficace, glaneur de cellule (GC)

* sans objets, foncteurs, labels, variants polymorphes, ...

```
let plateforme centre rayon alpha_init vitesse =  
  let alpha = ref alpha_init in  
  while true do  
    alpha := move !alpha;  
    draw centre rayon !alpha;  
  done
```

```
let plateforme centre rayon alpha_init vitesse =  
  let alpha = ref alpha_init in  
  while true do  
    alpha := move !alpha;  
    draw centre rayon !alpha;  
  done
```

```
let main =  
  Thread.create (plateforme c1 r a1) vitesse;  
  Thread.create (plateforme c2 r a2) vitesse
```

```
let plateforme centre rayon alpha_init vitesse =  
  let alpha = ref alpha_init in  
  while true do  
    alpha := move !alpha;  
    draw centre rayon !alpha;  
    Thread.yield()  
  done  
  
let main =  
  Thread.create (plateforme c1 r a1) vitesse;  
  Thread.create (plateforme c2 r a2) vitesse
```

```
let plateforme centre rayon alpha_init vitesse m1 m2 =  
  let alpha = ref alpha_init in  
  while true do  
    alpha := move !alpha;  
    draw centre rayon !alpha;  
    Mutex.unlock m2; Mutex.lock m1  
  done  
  
let main =  
  let m1, m2 = Mutex.create (), Mutex.create () in  
  Mutex.lock m1; Mutex.lock m2;  
  Thread.create (plateforme c1 r a1) vitesse m1 m2;  
  Thread.create (plateforme c2 r a2) vitesse m2 m1
```

Synchrone/**Asynchrone**

```
let barriere n =  
  let mutex, attente = Mutex.create (), Mutex.create () in  
  Mutex.lock attente;  
  let nb_att = ref 0 in  
  fun () ->  
    Mutex.lock mutex;  
    incr nb_att;  
    if !nb_att = n then begin  
      for i = 1 to n-1 do Mutex.unlock attente done;  
      nb_att := 0; Mutex.unlock mutex  
    end else begin  
      Mutex.unlock mutex; Mutex.lock attente  
    end  
end
```



```
let stop = barriere 3
```

```
let plateforme centre rayon alpha_init vitesse =  
  let alpha = ref alpha_init in  
  while true do  
    alpha := move !alpha;  
    draw centre rayon !alpha;  
    stop ()  
  done
```

```
let main =  
  Thread.create (plateforme c1 r a1) vitesse;  
  Thread.create (plateforme c2 r a2) vitesse;  
  Thread.create (plateforme c3 r a3) vitesse
```

```
let process plateforme centre rayon alpha_init vitesse =  
  let alpha = ref alpha_init in  
  while true do  
    alpha := move !alpha;  
    draw centre rayon !alpha;  
    pause  
  done
```

```
let process main =  
  run (plateforme c1 r a1 vitesse)  
  || run (plateforme c2 r a2 vitesse)  
  || run (plateforme c3 r a3 vitesse)
```

Le modèle réactif synchrone

Caractéristiques

- ▶ Instants logiques
- ▶ Composition parallèle synchrone
- ▶ Diffusion instantanée d'événements
- ▶ Création dynamique de processus

Origines

- ▶ Esterel [G. Berry & *al.* 1983]
- ▶ ReactiveC [F. Boussinot 1991]
- ▶ SL [F. Boussinot & R. de Simone 1996]

Autres langages :

- ▶ SugarCubes, Simple, Fair Threads, Loft, FunLoft, Lurc, S-pi, ...

ReactiveML

Présentation du langage

ReactiveML : les processus

Déclaration de processus :

▶ `let process <id> { <pattern> } = <expr>`

Expressions de base :

▶ coopération : `pause`

▶ exécution : `run <expr>`

Composition :

▶ séquentielle : `<expr> ; <expr>`

▶ parallèle : `<expr> || <expr>`

ReactiveML : les communications

Déclaration d'un signal :

▶ `signal <id> default <value> gather <function>`

Émission d'un signal :

▶ `emit <signal> <value>`

Statut d'un signal :

▶ attente : `await <signal> (patt) in <expr>`

▶ test de présence : `present <signal> then <expr> else <expr>`

Création dynamique de plates-formes

```
let process read_click click =  
  loop  
    if Graphics.button_down() then emit click (Graphics.mouse_pos());  
    pause  
  end  
val read_click : ((int * int) , 'a) event -> unit process
```

Création dynamique de plates-formes

```
let rec process add click =  
  await click (x,y) in  
  run (plateforme (float x, float y) 150. 0. vitesse)  
  ||  
  run (add click)  
val add : ('a, (int * int)) event -> unit process
```


ReactiveML : structures de contrôle

Préemption

- ▶ `do <expr> until <signal> done`
- ▶ `do <expr> until <signal> -> <expr> done`
- ▶ `do <expr> until <signal>(<patt>) -> <expr> done`

Suspension

- ▶ condition d'activation : `do <expr> when <signal> done`
- ▶ interrupteur : `control <expr> with <signal> done`

```
let process generate_new_plateforme click key new_plateforme =  
  loop  
    await click (p1) in  
    do  
      await click (p2) in  
        emit new_plateforme (p1, p2)  
    until key(Key_ESC) done  
  end
```

Programmation événementielle

```
class generate_new_plateforme = object(self)
  val mutable state = 0
  val mutable last_click = (0, 0)

  method on_click pos =
    match state with
    | 0 -> last_click <- pos;
          state <- 1
    | 1 -> emit new_plateforme (last_click, pos);
          state <- 0

  method on_key_down k =
    match k with
    | Key_ESC -> state <- 0
    | _ -> ()

end
```

ReactiveML

Reconfiguration dynamique

rmltop : le toplevel ReactiveML

Basé sur l'idée des Reactive Scripts [Boussinot & Hazard 96]

Utile pour :

- ▶ comprendre le modèle réactif
- ▶ faire des expériences de reconfiguration dynamique
- ▶ concevoir des systèmes réactifs

```
signal kill
```

```
val kill : (int, int list) event
```

```
let process killable p =
```

```
  let id = gen_id () in print_endline ("["^(string_of_int id)^"]");
```

```
  do run p
```

```
  until kill(ids) when List.mem id ids done
```

```
val killable : unit process -> unit process
```

Création dynamique : rappel

```
let rec process extend to_add =  
  await to_add(p) in  
  run p || run (extend to_add)  
val extend : ('a, 'b process) event -> unit process  
  
signal to_add  
  default process ()  
  gather (fun p q -> process (run p || run q))  
val add_to_me : (unit process, unit process) event
```

Création dynamique avec état

```
let rec process extend to_add state =  
  await to_add(p) in  
  run (p state) || run (extend to_add state)  
val extend : ('a , ('b -> 'c process)) event -> 'b -> unit process  
  
signal to_add  
  default (fun s -> process ())  
  gather (fun p q s -> process (run (p s) || run (q s)))  
val to_add : (('state -> unit process) , ('state -> unit process)) event
```


extensible

```
signal add
```

```
val add : ((int * (state -> unit process)),  
           (int * (state -> unit process)) list) event
```

```
let process extensible p_init state =
```

```
  let id = gen_id () in print_endline ("{"^(string_of_int id)^"}");
```

```
  signal add_to_me
```

```
    default (fun s -> process ())
```

```
    gather (fun p q s -> process (run (p s) || run (q s))) in
```

```
  run (p_init state) || run (extend add_to_me state)
```

```
  || loop
```

```
    await add(ids) in
```

```
    List.iter (fun (x,p) -> if x = id then emit add_to_me p) ids
```

```
  end
```

```
val extensible : (state -> 'a process) -> state -> unit process
```

ReactiveML

Parallélisme

ReactiveML et parallélisme

Exécution parallèle de programmes ReactiveML (Cédric Pasteur)

- ▶ simulation à grande échelle

Tâches asynchrones (Mehdi Dogguy)

- ▶ entrées/sorties asynchrones

Systèmes Globalement Asynchrones et Globalement Synchrones

- ▶ mélange ReactiveML et JoCaml

Conclusion

`http://rml.lri.fr`